Improving the Dynamic Creation of Processes in MPI-2

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This work has been partially supported by HP Brazil

IV Workshop PPD/UFRGS Porto Alegre, Brasil - Julho 2006

Outline

- Introduction
- 2 Programming with MPI-2
- On-line Scheduling in LAM-MPI
- Improving Process Scheduling in LAM-MPI
- 5 Conclusion

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 - Motivation
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Motivation

- MPI-2 Norm (1997)
 - Dynamic creation of processes
 - Remote Memory Access RMA
 - Parallel I/O
- The norm does not define any way to schedule dynamic processes
 - Which processor will receive a new process?
 - Inside a processor, in which order will the processes run?

Goal

- Offer on-line scheduling to MPI-2 programs
- Aiming at load balance
 - Number of processes created dynamically on each processor.

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MPI Distributions and MPI-2

- Some MPI distributions have started providing MPI-2 for a few years.
 - LAM-MPI, since the 2000's;
 - Also implements some tools to manage the dynamic entry of resources and their exit (lamgrow/lamshrink);
 - Based on daemons.
 - MPI-CH, since Jan., 2005.
 - HP-MPI, since Dec. 2005.
- MPI could almost be used on Grids:
 - It needs support for heterogeneity (MPI-G2);
 - It needs support for Fault-Tolerance (MPI-CHv2, MPI-FT).
- Open-MPI is a merge of LAM and MPI-FT, which could bring everything in a single distribution!

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MPI_Comm_spawn(char* command, char** argv,
int maxprocs, MPI_Info info, int root,
MPI_Comm comm, MPI_Comm *intercomm, int
*errcodes)

- MPI_Comm_spawn(char* command, char** argv, int maxprocs, MPI_Info info, int root, MPI_Comm comm, MPI_Comm *intercomm, int *errcodes)
 - Name of a MPI executable
 - A multi-process program with MPI_Init and MPI_Finalize
 - This executable will be the child program
 - Arguments of executable program
 - Command line parameters of the child program
 - Number of processes that will execute the child program

- MPI_Comm_spawn(char* command, char** argv, int maxprocs, MPI_Info info, int root, MPI_Comm comm, MPI_Comm *intercomm, int *errcodes)
 - Startup hints
 - For resource allocation LAM-MPI offers MPI_Info keys set by MPI_Info_set
 - lam_spawn_file defines a file appschema with available nodes
 - lam_spawn_sched_round_robin uses the LAM nodes making a Round-Robin distribution that starts on a determined node
 - MPI_INFO_NULL makes a Round-Robin distribution start on node with the lowest rank

MPI_Comm_spawn(char* command, char** argv, int maxprocs, MPI_Info info, int root, MPI_Comm comm, MPI_Comm *intercomm, int *errcodes)

- Parent intracommunicator
 - Children discover this intracommunicator by MPI_Comm_get_parent
- Child intercommunicator containing spawned processes
 - By this intracommunicator the parent exchanges messages with its children

The Fibonacci Example with MPI-2

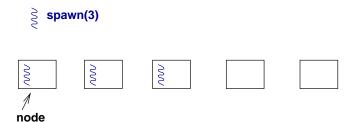
MPI-2 Code of the "fibo.c" program:

```
MPI Comm get parent(&parent);
if (n < 2) {
  MPI Isend (&n. 1, MPI LONG, 0, 1, parent, &reg);
else{
  sprintf (argv[0], "%ld", (n - 1));
  MPI Comm spawn ("Fibo", argv, 1, local info, myrank, MPI COMM SELF, &children comm[0],
                                                                        errcodes);
  sprintf (argv[0], "%ld", (n - 2));
  MPI Comm spawn ("Fibo", argy, 1, local info, myrank, MPI COMM SELF, &children comm[1],
                                                                        errcodes);
  MPI Recv (&x, 1, MPI LONG, MPI ANY SOURCE, 1, children comm[0], MPI STATUS IGNORE);
  MPI Recv (&y, 1, MPI LONG, MPI ANY SOURCE, 1, children comm[1], MPI STATUS IGNORE);
  fibn = x + y;
  MPI Isend (&fibn, 1, MPI LONG, 0, 1, parent, &reg);
MPI Finalize ();
```

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 - How to Schedule the Spawned Processes
- Improving Process Scheduling in LAM-MP
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- LAM-MPI provides a Round-Robin mechanism
 - It distributes maxprocs processes, one by node available
 - MPI_Info_set(info, "lam_spawn_sched_round_robin", node)



If there is a loop structure spawning processes

```
spawn(3)
spawn(3)
spawn(3)

A
node
```

- The distribution is not balanced needs to know who has received processes
- In a distributed case the problem is even bigger

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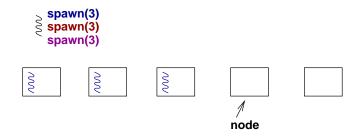
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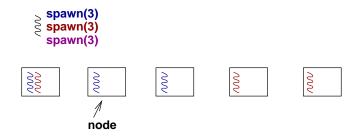
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Experimental Tests

 Spawning 20 processes on 5 nodes using single and multiple spawn calls with LAM Round-Robin mechanism

Environment	Node 1	Node 2	Node 3	Node 4	Node 5
20 spawns of 1 process	20	0	0	0	0
1 spawn of 20 processes	4	4	4	4	4

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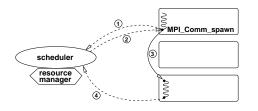
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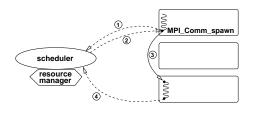
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 - Balancing the Load of MPI-2 Programs
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The Proposed Scheduler

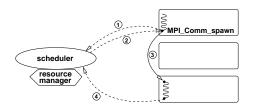
- A centralized deamon receives messages from re-defined MPI-2 primitives and takes the scheduling decisions
- The usage of pre-compilation redefinition is simple and portable.
 - the end-user just has to re-compile his program.



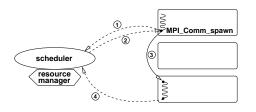
- Notification of processes creation
- Physical location is returned to parent
- The child is physically spawned
- Notification of processes completion



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Scheduling Heuristics

- Scheduling heuristics can be applied in two levels:
 - Scheduling processes into resources
 - Priorizing the execution of ready processes into a resource (actually it is left under OS responsibility)
- The scheduler implements our Round-Robin mechanism
 - The scheduler knows the last resource used
 - new_resource = (last_resource + 1)%total_resources
- Scheduling decisions according to the resources load information
 - Load Monitor collect the load information
 - Resource Manager coordinate of the load monitors and keeping a list of resources updated

 Comparing different schedules: number of processes spawned on each node

Environment	Node 1	Node 2	Node 3	Node 4	Node 5
fib(6) with LAM standard scheduler	25	0	0	0	0
fib(6) with embedded scheduler	8	4	8	2	3
fib(6) with proposed scheduler	5	5	5	5	5

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- Sets the first Round-Robin node to neighbor

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- Sets always the same first Round-Robin node
- Sets the first Round-Robin node to neighbor
- Sets the first Round-Robin node as global information

- LAM-MPI imposes a file description limitation
- Restrictions about the maximum number of processes running in a resource
- The Round-Robin mechanism proposed makes it possible to compute the 13th Fibonacci number

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	Node 1	Node 2	Node 3	Node 4	Node 5	Total Number of Processes
fib(13)	151	151	151	150	150	753

2nd Test-Case: an Irregular Computation

- Computation of prime numbers in a recursive search like Fibonacci program but irregular
- This program is CPU-intentive and a good load balance should impact the running time.
- Intervals range between 1 and 20 millions

Environment	Node 1	Node 2	Node 3	Node 4	Node 5
LAM standard scheduler	39	0	0	0	0
proposed scheduler	8	8	8	8	7

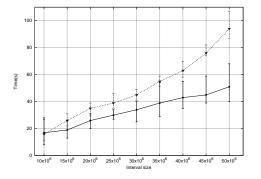
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LAM standard scheduler	39	0	0	0	0
proposed scheduler	8	8	8	8	7

- Good load balance with proposed scheduler
 - 181.15s with LAM standard scheduler
 - 46.12s with proposed scheduler

List-scheduleing – Using Dynamic Load Information



 Round-Robin strategy (dotted line) and List scheduling (solid line)

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Conclusion

- This work aims to simplify the on-line scheduling of MPI-2 programs
 - Interest for dynamic platforms!
- The native LAM implementation is not efficient, due to a simple scheduling strategy.
 - A simple prototype led to clear improvements

Next Steps

- To implement the proposed scheduler inside of LAM-MPI distribution
- Tests with real-world applications
 - Branch & Bound, linear systems solving, etc...
- To distribute the centralized scheduler
- To implement a workstealing strategy for the distributed scheduler

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